

# Characterising Continuously Realizable Multifunctions

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## Survey

- ▶ Multivalued functions
- ▶ The Characterisation Theorem by Brattka and Hertling
- ▶ The powerspace  $\mathcal{K}_{+b}(Y)$
- ▶ A more general Characterisation Theorem

# Multifunctions

## Remember

The following maps are incomputable:

- ▶ Tests like  $x < y$ ,  $x \geq y$  etc.
- ▶ Zero-finding for polynomials:
  - ▶ There is no computable selection function  $s: \mathbb{R}^3 \rightarrow \mathbb{R}$  s.t.  $s(a, b, c)$  is a zero of  $x^3 + ax^2 + bx + c$ .

## Alternative

- ▶ Replace  $x < y$  by the *finite precision test* ( $x <_\varepsilon y$ ):

$$(x <_\varepsilon y) := \begin{cases} \{\text{true}\} & \text{if } x \leq y - \varepsilon \\ \{\text{true}, \text{false}\} & \text{if } y - \varepsilon < x < y \\ \{\text{false}\} & \text{if } x \geq y \end{cases}$$

- ▶ Compute the relation  $Z \subseteq \mathbb{R}^3 \times \mathbb{R}$  given by

$$(a, b, c) Z x \quad \text{iff} \quad x^3 + ax^2 + bx + c = 0$$

## Definition

- ▶ A *multifunction* (or *computational task*)  $F$  is a relation between represented spaces  $X$  and  $Y$ , written as  $F: X \rightrightarrows Y$ .
- ▶  $X$  is the *input space*,  $Y$  is the *output space* of  $F$ .
- ▶ Notation:  $F[x] := \{y \in Y \mid (x, y) \in F\}$ .
- ▶  $H: X \rightrightarrows Y$  is a *tightening* for  $F$ , if  $\forall x \in X. \emptyset \neq H[x] \subseteq F[x]$ .

## Remark

We will assume every multifunction to be *total*, i.e.  $F[x] \neq \emptyset$  for all  $x \in X$ .

## Remember

A *represented space*  $X$  is a set endowed with a representation  $\delta_X: \mathbb{N}^{\mathbb{N}} \dashrightarrow X$ .

### Definition

Let  $F: X \rightrightarrows Y$  be a total multifunction between represented spaces  $X, Y$ .

- ▶  $g: \mathbb{N}^{\mathbb{N}} \dashrightarrow \mathbb{N}^{\mathbb{N}}$  is a *realizer* for  $F$ , if

$$\delta_Y g(p) \in F[\delta_X(p)] \quad \text{for all } p \in \text{dom}(\delta_X).$$

- ▶ Diagram:

$$\begin{array}{ccc}
 X & \xRightarrow{F} & Y \\
 \delta_X \uparrow & \circlearrowleft & \uparrow \delta_Y \\
 \text{dom}(\delta_X) & \xrightarrow{g} & \text{dom}(\delta_Y)
 \end{array}$$

- ▶  $F$  is called *computable (continuously realizable)*, if  $F$  has a computable (continuous) realizer  $g$ .

## Observation

Any realizer  $g$  for a tightening  $H$  for  $F$  is a realizer for  $F$ .

## Proof

Since  $\delta_Y g(p) \in H[\delta_X(p)] \subseteq F[\delta_X(p)]$ .

## Lemma

- ▶ Computationally equivalent representations induce the same notion of computable multifunctions.
- ▶ Topologically equivalent representations induce the same notion of continuously realizable multifunctions.

**Example** (Computable multifunctions)

- ▶ The finite precision test  $(x <_\varepsilon y): \mathbb{R} \times \mathbb{R}_{>0} \times \mathbb{R} \rightrightarrows \mathbf{Bool}$ ,

$$(x <_\varepsilon y) \ni \begin{cases} \text{true} & \text{if } x < y \\ \text{false} & \text{if } x > y - \varepsilon \end{cases}$$

- ▶ Zero-finding  $Z_3: \mathbb{R}^3 \rightrightarrows \mathbb{R}$  of polynomials of degree 3,

$$Z_3[a, b, c] := \{x \in \mathbb{R} \mid x^3 + ax^2 + bx + c = 0\}$$

- ▶ The inverse  $\delta_X^{-1}: X \rightrightarrows \text{dom}(\delta_X)$  of the representation  $\delta_X$ ,

$$\delta_X^{-1}[x] := \{p \in \mathbb{N}^{\mathbb{N}} \mid \delta_X(p) = x\}$$



## Main Question

Find a characterisation of continuously realizable multifunctions in terms of “mathematical” continuity.

*Remember:*

## Theorem

Let  $X, Y$  be *admissibly* represented spaces.

Then a function  $f: X \rightarrow Y$  has a continuous realizer iff  $f$  is sequentially continuous.

## Multifunctions with compact images

## Definition

A total multifunction  $H: X \rightrightarrows Y$  has *compact images*, if  $H[x]$  is compact in  $Y$  for all  $x \in X$ .

## Proposition

Let  $X$  be a computable metric space.

Any computable  $F: X \rightrightarrows Y$  has a computable tightening  $H: X \rightrightarrows Y$  with compact images.

*The proof is based on:*

## Proposition

Any computable metric space  $X$  has a representation with compact fibers which is computably equivalent to its Cauchy representation.

## Observation

A total multifunction  $H: X \rightrightarrows Y$  with compact images can be identified with the function  $h: X \rightarrow \mathbf{K}(Y)$ ,  $x \mapsto H[x]$ .

## Definition

- ▶  $\mathbf{K}(Y) := \{\text{all non-empty compact subsets of } Y\}$
- ▶ Basic sets of the *lower Vietoris topology* on  $\mathbf{K}(Y)$ :

$$\diamond V := \{K \text{ compact} \mid K \cap V \neq \emptyset\}$$

where  $V$  is open in  $Y$ .

- ▶  $h: X \rightarrow \mathbf{K}(Y)$  is *lower semi-continuous*, if  $h$  is continuous w.r.t. the lower Vietoris topology.

**Characterisation Theorem** (V. Brattka & P. Hertling 1994)

Let  $X, Y$  be separable metric spaces.

Then  $F: X \rightrightarrows Y$  has a continuous realizer if, and only if,  $F$  has a lower semi-continuous tightening  $h: X \rightarrow \mathbb{K}(Y)$ .

## Proposition

- ▶ Let  $Y$  be a  $\text{QCB}_2$ -space without a countable base.
- ▶ Let  $\kappa_+$  be the positive representation for  $\mathbf{K}(Y)$ .

Then

- ▶  $\kappa_+$  is lower semi-continuous.
- ▶ But viewed as a multifunction  $\kappa_+ : \text{dom}(\kappa_+) \rightrightarrows Y$ ,  $\kappa_+$  has no continuous realizer.

## Remark

- ▶ This means that compact choice for  $Y$  is not continuously realizable w.r.t.  $\kappa_+$ .
- ▶ *Compact choice*: selecting an element in a compact set.
- ▶ By contrast, compact choice for any computable metric space  $M$  is computable w.r.t. its  $\kappa_+$ .

How can we handle output spaces  $Y \notin \omega\text{Top}$ ?

**The powerspace**  $\mathcal{K}_{+b}$

## Aim

Construct a more informative representation for  $\mathbf{K}(Y)$  admitting continuous realizability of compact choice.

## Definition

Let  $Y$  be an admissibly represented  $\mathbf{QCB}_2$ -space.

- ▶ Define a representation  $\kappa_{+b}$  for  $\mathbf{K}(Y)$  by

$$\kappa_{+b}\langle p, b \rangle = K \quad \text{iff} \quad \kappa_+(p) = K \ \& \ K \subseteq \kappa_-(b)$$

where  $\kappa_+, \kappa_-$  are the canonical positive / negative representations for  $\mathbf{K}(Y)$  used in Computable Analysis.

- ▶ Define  $\mathcal{K}_{+b}(Y) := (\mathbf{K}(Y), \kappa_{+b})$ .
- ▶ Define  $\mathcal{K}_+(Y) := (\mathbf{K}(Y), \kappa_+)$ .



**Lemma** (The convergence relation)

A sequence  $(K_n)_n$  converges to  $K_\infty$  in  $\mathcal{K}_{+b}(Y)$  iff:

- (a)  $(K_n)_n$  converges to  $K_\infty$  w.r.t. the lower Vietoris topology &
- (b)  $\bigcup_{n \in \mathbb{N}} K_n$  is relatively compact.

**Example**

- ▶  $([0, n])_n$  converges to  $[0, 1]$  in  $\mathcal{K}_+(\mathbb{R})$ ,
- ▶ but it is divergent in  $\mathcal{K}_{+b}(\mathbb{R})$ .
- ▶  $([0, 1 - 2^{-n}])_n$  converges to  $[0, 1]$  in  $\mathcal{K}_{+b}(\mathbb{R})$ ,
- ▶ but not to  $[0, 2]$ .

## Quasi-normal spaces (CCA 2008)

- ▶ Quasi-normal space = a QCB-space that arises as the sequentialisation of a normal space.
- ▶ Examples:
  - ▶ All separable metric spaces
  - ▶ The space  $\mathcal{D}$  of distributions
  - ▶ The Kleene-Kreisel spaces  $\mathbb{N}^{\mathbb{N}^{\mathbb{N}}}, \mathbb{N}^{\mathbb{N}^{\mathbb{N}^{\mathbb{N}}}}, \dots$
  - ▶ All Co-Polish spaces
- ▶ Quasi-normal spaces have excellent closure properties:
  - ▶ cartesian closed
  - ▶ countable products and equalisers ( $\approx$  subspaces)
  - ▶ countable co-products and co-equalisers
- ▶ Quasi-normal spaces comprise many Hausdorff spaces used in Computable Functional Analysis.

## Proposition

Let  $Y$  be quasi-normal.

- ▶  $\mathcal{K}_{+b}(Y)$  is an admissibly represented limit space.
- ▶  $\mathcal{K}_{+b}(Y)$  is not topological, unless  $Y$  is compact.
- ▶ Compact choice for  $Y$  is continuously realizable w.r.t.  $\kappa_{+b}$  (i.e. there is a continuous selector  $S: \text{dom}(\kappa_{+b}) \rightarrow Y$  such that  $S(\rho) \in \kappa_{+b}(\rho)$ ).

## Proposition

Let  $X$  have a countable base, let  $Y$  be quasi-normal.

Every continuous function  $f: X \rightarrow \mathcal{K}_{+b}(Y)$  can be ‘blown up’ to a upper semi-continuous function  $\bar{f}: X \rightarrow \mathcal{K}_-(Y)$  such that  $\forall x \in X. f(x) \subseteq \bar{f}(x)$ .

## Characterisation Theorem

Let  $X$  be separable metrisable, let  $Y$  be quasi-normal.

Then  $F: X \rightrightarrows Y$  has a continuous realizer if, and only if,  
 $F$  has a sequentially continuous tightening  $h: X \rightarrow \mathcal{K}_{+b}(Y)$ .

## Remark

- ▶ “ $\Leftarrow$ ” holds for any represented space  $X$ .
- ▶ “ $\Rightarrow$ ” holds for any admissibly represented  $\text{QCB}_2$ -space  $Y$ .

## Question

Can we generalise this to input spaces  $X \notin \omega\text{Top}$  ?

## Counterexample

Let  $\text{Poly}$  be the Co-Polish space of polynomials.

Define  $\text{DegreeBound}: \text{Poly} \rightrightarrows \mathbb{N}$  by

$$\text{DegreeBound}[P] \ni k : \iff k \geq \text{Degree}(P).$$

Then

- ▶  $\text{DegreeBound}$  is computable.
- ▶ But  $\text{DegreeBound}$  does not have a continuously realizable tightening with compact images.

## Lemma

For every non-metrisable Co-Polish space  $X$  there is

- ▶ a computable  $F: X \rightrightarrows \mathbb{N}$
- ▶ that does not have a continuously realizable tightening with compact images.

## Remember

A *Co-Polish space* is the direct limit of an increasing sequence of compact metrisable subspaces  $X_k$ .

## Lemma

For every  $X \in \text{QCB}_1 \setminus \omega\text{Top}$  there is

- ▶ a continuously realizable  $F: X \rightrightarrows \mathbb{N}$
- ▶ that does not have a lower semi-continuous tightening  $h: X \rightarrow \mathbb{K}(\mathbb{N})$ .

## Simple Fact

The inverse  $\delta_X^{-1} : X \rightrightarrows \text{dom}(\delta_X)$  is computable.

## Proposition

Let  $X$  be an admissibly represented  $\text{QCB}_2$ -space. TFAE:

- (a)  $\delta_X^{-1}$  has a continuous tightening  $h : X \rightarrow \mathcal{K}_{+b}(\text{dom}(\delta_X))$ .
- (b)  $X$  is separable metrisable.

## Corollary

The ‘only-if-part’ of the Characterisation Theorem holds for metrisable input spaces  $X$  *only*.

## Summary

- ▶ Multivalued functions form an indispensable tool in Computable Analysis.
- ▶ A computable multifunctions on a computable metric space  $X$  has a computable tightening with compact images.
- ▶ The powerspace  $\mathcal{K}_{+b}(Y)$  has nice properties (at least) for quasi-normal spaces.
- ▶  $\mathcal{K}_{+b}(Y)$  allows us to characterise continuously realizable multifunctions from metrisable spaces to quasi-normal spaces.
- ▶ Open problem:  
Find a characterisation of continuously realizable multifunctions on non-metrisable input spaces.