

# Computable Multifunctions on Effectively Hausdorff Spaces

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We continue our research on characterising realizable multifunctions (cf. [5]). Computable multifunctions play a prominent rôle in Computable Analysis [1, 7]. They are used as an appropriate substitute for single-valued functions that fail to be computable just for topological reasons. For example, the discontinuous and thus incomputable test “ $x < y$ ” on the real numbers can often be replaced by the computable *finite precision test*  $(x <_\varepsilon y) : \mathbb{R} \times \mathbb{R}_{>0} \times \mathbb{R} \rightrightarrows \{\mathbf{tt}, \mathbf{ff}\}$  which allows both answers for inputs satisfying  $y - \varepsilon < x < y$ . (cf. [1]).

We introduce a new notion of an effectively Hausdorff space which is stronger than the current notion of a computable Hausdorff space proposed by A. Pauly in [3]. We then characterise computable multifunctions from computable metric spaces to effectively Hausdorff spaces. Inspired by this characterisation, we propose a stronger computability notion for multifunctions than the usual one and study some of its properties.

## The usual notion of a computable multifunction

A multifunction  $F$  between admissibly represented spaces  $\mathsf{X}, \mathsf{Y}$  is a binary relation between the underlying sets of  $\mathsf{X}$  and  $\mathsf{Y}$ , written as  $F : \mathsf{X} \rightrightarrows \mathsf{Y}$ . The elements of  $F[x] := \{y \in \mathsf{Y} \mid (x, y) \in F\}$  are viewed as the legitimate results for the input  $x$ . A multifunction  $F$  is called *computable* [6, 7], if there is a computable function  $g : \mathbb{N}^{\mathbb{N}} \rightarrow \mathbb{N}^{\mathbb{N}}$  satisfying

$$\delta_{\mathsf{Y}}(g(p)) \in F[\delta_{\mathsf{X}}(p)] \quad \text{for every } p \in \text{dom}(F\delta_{\mathsf{X}}).$$

Here  $\delta_{\mathsf{X}}$  and  $\delta_{\mathsf{Y}}$  denote the respective representations of  $\mathsf{X}$  and  $\mathsf{Y}$ .

## Effectively Hausdorff spaces

It is known that the topology of every Hausdorff qcb-space contains a countably-based Hausdorff subtopology. We employ this fact in our definition of an effectively Hausdorff space.

**Definition 1** We call a represented space  $\mathsf{X}$  an *effectively Hausdorff space*, if there are two computable sequences  $(u_i)_i, (v_i)_i$  of open subsets of  $\mathsf{X}$  such that  $\bigcup_i (u_i \times v_i) = \mathsf{X}^2 \setminus \{(x, x) \mid x \in \mathsf{X}\}$  and the representation  $\delta_{\mathsf{X}}$  is computably admissible.

Any computable metric space is an example of an effectively Hausdorff space. Effectively Hausdorff spaces are computable Hausdorff spaces in the sense of [3], but not vice versa. Indeed, every effectively Hausdorff space is topologically Hausdorff, whereas computable Hausdorffness only entails sequential Hausdorffness (meaning that every convergent sequence has a unique limit). Effective Hausdorffness admits effectivisations of classical theorems from topology: for example the fact that compact Hausdorff spaces are regular can be effectivised by stating that any computably compact, effectively Hausdorff space is computably regular<sup>1</sup>.

We apply effective Hausdorffness in the following characterisation theorem of computable multifunctions.

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<sup>1</sup>See [3] for the definition of computable compactness and [4] for computable regularity.

**Theorem 2** *Let  $X$  be a computable metric space, and let  $Y$  be an effectively Hausdorff space. Then a total multifunction  $F: X \rightrightarrows X$  is computable iff there are two computable functions  $h_+: X \rightarrow \mathcal{A}_+(Y)$  and  $h_b: X \rightarrow \mathcal{K}_-(Y)$  satisfying  $\emptyset \neq h_+(x) \subseteq F[x] \cap h_b(x)$  for all  $x \in X$ .*

Here  $\mathcal{A}_+(Y)$  denotes the represented space of non-empty closed subsets of  $Y$  with the positive (overt) representation and  $\mathcal{K}_-(Y)$  the represented space of compact subsets of  $Y$  with the negative representation (cf. [2, 3, 6]). Note that  $h_+(x)$  is in fact compact by being a closed subset of a compact set.

### Strongly computable multifunctions

The only-if-part of Theorem 2 only holds for metrisable spaces. We now introduce a refined notion of computability for multifunctions which is based on computability of set-valued functions and works appropriately at least for multifunctions between effectively Hausdorff spaces.

**Definition 3** We call a total multifunction  $F$  *strongly computable*, if there are two computable functions  $h_+: X \rightarrow \mathcal{A}_+(Y)$ ,  $h_-: X \rightarrow \mathcal{K}_-(Y)$  satisfying  $\emptyset \neq h_+(x) \subseteq F[x] \subseteq h_-(x)$  for all  $x \in X$ .

This computability notion chimes with a composition  $\diamond$  for multifunctions defined by

$$G \diamond F[x] := \text{Cls} \bigcup \{G[y] \mid y \in F[x]\}.$$

If the involved spaces are effectively Hausdorff, then  $\diamond$  preserves strong computability; moreover, any strongly computable multifunction is computable in the usual sense and computable total functions form strongly computable multifunctions. We obtain:

**Proposition 4** *Effectively Hausdorff spaces together with strongly computable multifunctions having compact images form a category (which we denote by  $\text{EffHausMulti}$ ). It has equalisers and finite coproducts.*

### Discussion

Any (in the usual sense) computable multifunction from a computable metric space  $X$  to an effectively Hausdorff space has a tightening which is strongly computable. This is not true, if  $X$  is not metrisable. The category  $\text{EffHausMulti}$  is not cartesian, it only has ‘weak’ products. It is not known whether effective Hausdorffness is necessary for Theorem 2 or Proposition 4.

## References

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