#### A Verified Power-Series Method for Multivariate IVPs

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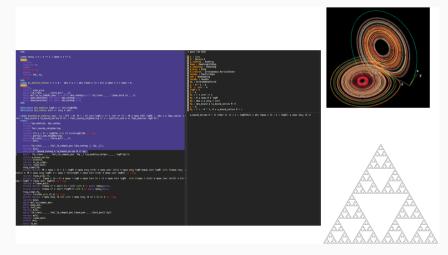






## **Computation and Formal Proof**

A **proof assistant** is an <u>interactive</u> tool for writing and checking formal proofs.



### **Ordinary Differential Equations (ODEs)**

In this talk we consider initial value problems (IVPs) for ordinary differential equations (ODEs). An IVP is an ODE together with an initial condition at

$$\dot{\vec{y}}(t) = \vec{F}(t, \vec{y}(t)), \quad \vec{y}(t_0) = \vec{y}_0, \quad \vec{F}: D \subseteq \mathbb{R}^{d+1} \to \mathbb{R}^d.$$

A solution is a differentiable function  $\vec{y}:I\to\mathbb{R}^d$ , where  $I\subseteq\mathbb{R}$  is an open interval containing  $t_0$ , that satisfies the IVP.

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#### **Remarks:**

- Higher-order ODEs  $\vec{y}^{(k)} = \vec{G}(t, \vec{y}, \dots, \vec{y}^{(k-1)})$  can be rewritten as a first-order ODEs by introducing additional variables for the higher derivatives  $\vec{y}^{(2)}, \dots, \vec{y}^{(k)}$ .
- ullet Any ODE can be turned into an equivalent autonomous ODE by adding t as additional variable.
- By translation, one can assume  $t_0=0$  and  $\vec{y}_0=\vec{0}$  without loss of generality.

## **ODE solving in Proof Assistants**

The problem is a natural candidate for formal verification:

- It is a classical topic in analysis with many applications in science and engineering.
- Most ODEs do not have a closed-form solution, therefore numerical methods are required to approximate solutions.
- Guaranteeing correctness of these approximations is important for safety-critical applications.

Although other proof assistants (e.g. Isabelle/HOL) include substantial formalizations of ODE theory and numerical methods, not so much is available in Rocq so far.

• The CoRN library includes an implementation of the Picard Iteration method. [1]

<sup>[1]</sup> Makarov & Spitters. The Picard algorithm for ordinary differential equations in Coq, Proc. of ITP 2013.

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Since Rocq is based on constructive logic, a constructive/computable formalization is natural.

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### **Computability and Complexity of ODEs**

Assume 
$$f:[0,1]\times[-1,1]\to\mathbb{R}$$
 and consider the IVP  $\dot{y}(t)=f(t,y(t));y(0)=0$ 

Assumptions on f	Computability/Complexity of Solution
(Polynomial-time) computable	All solutions can be non-computable [3]
+ Unique solution	Computable, but complexity can be arbitrarily high [4]
+ Lipschitz continuous	PSPACE complete [5]
+ Analytic	Solution is polynomial-time computable [6]

<sup>[3]</sup> Aberth, The failure in computable analysis of a classical existence theorem for differential equations, 1971.

 $\hbox{ [6] M\"{u}ller \& Moiske, Solving initial value problems in polynomial time, Proc. of JAIIO-Panel '93, 1993. }$ 

<sup>[4]</sup> Ko, On the computational complexity of ordinary differential equations, Journal of Computer and System Sciences, 1983.

<sup>[5]</sup> Kawamura, Lipschitz Continuous Ordinary Differential Equations are Polynomial-Space Complete, Comp. Complexity, 2010.

#### The Cauchy-Kovalevskaya Theorem

#### Theorem (ODE version of the Cauchy-Kovalevskaya Theorem)

Let  $U\subseteq\mathbb{R}^d$  be open and  $f:U\to\mathbb{R}^d$  analytic. Then the initial value problem

$$\dot{\vec{y}}(t) = \vec{F}(\vec{y}(t)), \quad \vec{y}(0) = \vec{y}_0$$

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Main idea of the classical proof:

- Expand  $\vec{F}(\vec{y})$  in a formal power series around  $\vec{y_0}$ .
- Derive a recurrence for the Taylor coefficients of  $\vec{y}(t)$  from the ODE  $\dot{\vec{y}}=\vec{F}(\vec{y})$ .
- Prove that the formal series has a positive radius of convergence by dominating it by a geometric <u>majorant sequence</u>.

#### **Formal Power Series Solution**

Let

$$\dot{\vec{y}}(t) = \vec{F}(\vec{y}(t)), \quad \vec{y}(0) = \vec{y}_0.$$

We inductively define a sequence of functions  $\vec{F}^{[n]}:\mathbb{R}^d o \mathbb{R}^d$  of multivariate functions by

$$ec{F}^{[0]}(ec{x})=ec{x}$$
 , and 
$$ec{F}^{[n+1]}(ec{x})=J_{ec{F}^{[n]}}(ec{x})\cdotec{f}(ec{x})$$

Define a function by the formal power series

$$\vec{y}(t) = \sum_{i=0}^{\infty} \frac{1}{n!} \vec{F}^{[n]}(\vec{y_0}) t^n$$

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• A (single-variate) power series  $M(x)=\sum_{k=0}^\infty A_k x^k$  is said to majorize a d-variate power series  $f(x)=\sum_{\alpha\in\mathbb{N}^d}a_\alpha x^\alpha$  if

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- In that case, the tail of the series satisfies

$$\left| \sum_{|\alpha| > k} a_{\alpha} x^{\alpha} \right| \leq M \sum_{n=k+1}^{\infty} (R||x||)^{n} = M \frac{(R||x||)^{k+1}}{1 - R||x||}, \quad (R||x|| < 1).$$

• We can show that if the right-hand side function  $f:\mathbb{R}^d\to\mathbb{R}^d$  has a geometric majorant (M,R) then (1,2dMR) works as a geometric majorant for the solution y, guaranteeing convergence of y(t) for all t with  $|t|<\frac{1}{2dMR}$ .

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- We get radius  $\frac{1}{4}$  for the solution, while the actual solution  $y(t) = \frac{1}{1-t}$  has radius 1.
- On the other hand for

$$\dot{y}(t) = \frac{1}{1 - y(t)} = \sum_{k=0}^{\infty} y(t)^k \; ; \; y(0) = 0$$

we get radius  $\frac{1}{2}$  which is identical to the actual radius of the solution

$$y(t) = 1 - \sqrt{1 - 2t}.$$

#### **Polynomial IVPs**

- ODEs with polynomial right-hand side are an important special case.
- Polynomials allow efficient implementations of evaluation and other operations.
- As the coefficient sequence is finite, majorant bounds can be computed automatically.
- Many analytic ODEs (including all systems built from elementary functions) can be rewritten as polynomial ODEs by increasing the dimension.

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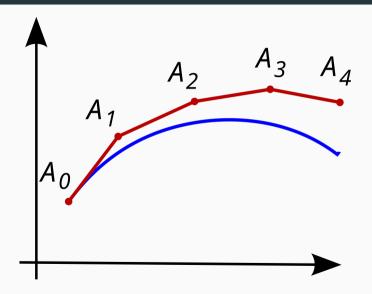
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**Example:**  $y_1' = \frac{1}{2y_1}$ ;  $y_1(0) = 1$  The IVP has the solution  $y_1(x) = \sqrt{x+1}$ .

Introduce auxiliary variable  $y_2 = \frac{1}{y_1}$  to obtain:

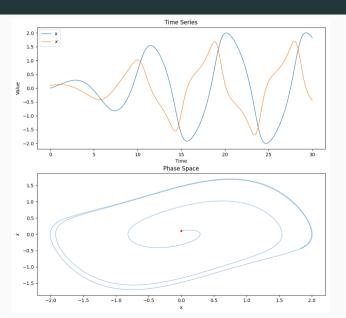
$$\dot{y_1} = \frac{1}{2} y_2,$$
  $y_1(0) = 1,$   $\dot{y_2} = -\frac{1}{2} y_2^3,$   $y_2(0) = 1.$ 

# **Extending the solution**





#### Demo



• Formalized a constructive version of the Cauchy-Kovalevskiya theorem in Rocq.

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- Allows to compute interval trajectories inside the proof assistant.
- Currently only very crude error bounds, sharpening would improve efficiency.
- In general, optimizing computation inside the Rocq proof assistant is challenging.

# Thank you!



